CGAL - the Computational Geometry Algorithms Library
Pierre Alliez, Andreas Fabri, Efi Fogel

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Computational Geometry Algorithms Library

Pierre Alliez
INRIA

Andreas Fabri
GeometryFactory

Efi Fogel
Tel-Aviv University
Abstract

The CGAL Open Source Project provides easy access to efficient and reliable geometric algorithms in the form of a C++ library, offering geometric data structures and algorithms, which are efficient, robust, easy to use, and easy to integrate in existing software. The usage of de facto standard libraries increases productivity, as it allows software developers to focus on the application layer. This course is targeted at software developers with geometric needs, and course graduates will be able to select and use the appropriate algorithms and data structures provided by CGAL in their current or upcoming projects.

Key Facts

CGAL 3.3, released in June 2007: 90 software components, 600,000 lines of code, 3,500 user and reference manual pages, Cross platform support, Annual release with 12,000 downloads, 1,000 subscribers on the user mailing list, 40 subscribers on the developer mailing list. CGAL is used in many application areas by companies as Total (Oil&gas), British Telecom (Telecom), Cadence (VLSI), Leica Geosystems (GIS), Dassault Systèmes (CAD), The Moving Picture Company (Visual effects).

CGAL Project

The project is steered by an Editorial Board, it has a well defined development process, and the infrastructure for distributed development. The following research institutes and companies are actively involved or made contributions to the library: INRIA-Sophia-Antipolis, Max-Planck Institute for Computer Science, Tel-Aviv University, GeometryFactory, ETH Zurich, FU Berlin, University of Groningen, University of Utrecht, Stanford University, Athens University, and the Foundation of Research and Technology – Hellas. For more information on the project see www.cgal.org For an overview on what is in CGAL: http://www.cgal.org/Manual/3.3/doc_html/cgal_manual/packages.html
Organization of the Course

This course starts with an overview followed by three in-depth sessions covering central data structures: The overview session presents the CGAL project and the design principles of CGAL. CGAL has adopted the exact computing paradigm, which yields robust and at the same time fast algorithms. CGAL further has adopted the generic programming paradigm, which makes CGAL particularly easy to customize and to integrate. Finally, we show how CGAL fits naturally with the STL, and the Boost graph library.

The session on polyhedral surfaces presents the underlying halfedge data structure and how it can be customized to user needs. We further present algorithms for polyhedral surfaces like parameterization, mesh subdivision and simplification, Boolean operations, and intersection detection.

The session on arrangements presents the arrangement API and several data structures built on top of it. These are 3D Minkowski sums, which can be used for collision detection, and 3D lower envelopes, which can be used for visibility map computations.

The last session covers the 2D and 3D triangulation API as well as the surface and volume mesh generators, which are based on Delaunay refinement.

The source code of the examples will be made available at http://www.cgal.org/Courses/SIGGRAPH2008.
Presenters

Dr. Andreas Fabri (organizer)
Chief Officer
GeometryFactory
20, Avenue Yves Emmanuel Baudoin 06130 Grasse, FRANCE
Email: andreas.fabri@geometryfactory.com
http://www.geometryfactory.com/

Dr. Efi Fogel
Senior Researcher
Faculty of Exact Sciences. Tel Aviv University Tel Aviv 69978, ISRAEL
Email: efifogel@gmail.com
http://www.cs.tau.ac.il/~efif/

Dr. Pierre Alliez
Senior Researcher
INRIA Sophia Antipolis – Mediterranee
2004 route des Lucioles BP 93 - GEOMETRICA, FRANCE
Email: pierre.alliez@sophia.inria.fr
http://www-sop.inria.fr/geometrica/team/Pierre.Alliez/
**Biographies**

**Andreas Fabri**, PhD, GeometryFactory As member of the initial development team of the CGAL project, Andreas Fabri is one of the architects of the CGAL software. For several years he chaired the CGAL Editorial Board. In 2003 Andreas founded the GeometryFactory as spin-off of the CGAL project, offering licenses, service and support to commercial users, who cannot comply with the Open Source license of CGAL. Andreas received his PhD in computer science in 2004 from Ecole de Mines de Paris while working on geometric algorithms for parallel machines at INRIA.

**Pierre Alliez** obtained his PhD from Ecole nationale supérieure des Télécommunications, did his postdoc at Caltech, and is researcher at INRIA since 2001. His main research interests are on topics commonly referred to as Geometry Processing: geometry compression, surface approximation, mesh parameterization, surface remeshing and mesh generation. He is this year co-chair of the EUROGRAPHICS Symposium on Geometry Processing. In 2005 Pierre Alliez received the Eurographics Young Researcher Award.

**Efi Fogel**, MSc, Tel-Aviv University Efi Fogel is a co-founder of LucidLogix Ltd., a startup company that intends to deliver high performance 3D graphics systems. Efi Fogel is completing his Ph.D. studies at Tel-Aviv University. 3D Graphics and Computational Geometry are his main areas of interests. He is a member of the CGAL Editorial Board, and he is deeply involved with the design and implementation of the arrangement package of CGAL and its derivatives. Efi Fogel received his M.Sc. from Stanford University in 1989. He worked for Silicon Graphics Inc. (SGI) between 1989-1997 at the Advanced Graphics Devision, where he contributed to the specification of OpenGL among the other. After that Efi worked for Immersia Ltd, and he served as the CTO of Enbaya Ltd.
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Bibliography

- Bibliographic entries for individual chapters of CGAL manuals
- CGAL User and Reference Manual Bibliography
CGAL Contributors

CGAL Editorial Board
Sylvain Pion, Pierre Alliez, Eric Berberich, Andreas Fabri, Efi Fogel, Bernd Gärtner, Michael Hemmer, Michael Hoffmann, Menelaos Karavelas, Marc Pouget, Monique Teillaud, Ron Wein

CGAL Developers
Fernando Cacciola, Frédéric Cazals, Raphaëlle Chaine, Pavel Emeliyanenko, Marc Glisse, Luc Habert, Peter Hachenberger, Idit Haran, Thomas Herrmann, Samuel Hornus, Michael Kerber, Nico Kruithof, Sebastian Limbach, Sébastien Loriot, Pedro Machado Manhães de Castro, Andreas Meyer, Michał Meyerovitch, Thanh-Trung Nguyen, Luis Peñaranda, Joachim Reichel, Laurent Rineau, Daniel Russel, Laurent Sabore, Ophir Setter, Tel-Aviv University, George Tzoumas

Alumni

EU Project Board Members
Helmut Alt, Jean-Daniel Boissonnat, Dan Halperin, Kurt Mehlhorn, Stefan Näh, Mark Overmars, Geert Vegter, Emo Welzl, Peter Widmayer
Computational Geometry Algorithms Library

Pierre Alliez  Andreas Fabri  Efi Fogel
INRIA       GeometryFactory  Tel Aviv University
## Course Outline

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Overview

Andreas Fabri
GeometryFactory
Mission Statement

“Make the large body of geometric algorithms developed in the field of computational geometry available for industrial applications”

CGAL Project Proposal, 1996
Algorithms and Datastructures

Bounding Volumes

Polyhedral Surface

Boolean Operations

Triangulations

Voronoi Diagrams

Mesh Generation

Subdivision

Simplification

Parametrisation

Streamlines

Ridge Detection

Neighbor Search

Kinetic Datastructures

Lower Envelope

Arrangement

Intersection Detection

Minkowski Sum

PCA

Polytope distance

QP Solver
CGAL in Numbers

500,000 lines of C++ code
10,000 downloads/year (+ Linux distributions)
   3,500 manual pages
3,000 subscribers to cgal-announce
1,000 subscribers to cgal-discuss
   120 packages
      60 commercial users
   20 active developers
12 months release cycle
   2 licenses: Open Source and commercial
Some Commercial Users

- Cadence
- Pulsic
- Toshiba
- Orbotech
- IDE
- TOSYS
- QinetiQ
- BAE Systems
- Leica Geosystems
- SAFE Software
- Video Arch
- Safe Software
- Inpho
- Agilent Technologies
- VDRC
- Noesis
- ProtoMold
- Exa Corporation
- ZW CAD
- ECL
- MPC
- Dassault Systems
- Schaefer Mayfield
- Tomtec
- St. Jude Medical
- The MathWorks
- TGS
- TruePosition
- BT
- IFP
- Total
- BSAP
- GIS
- VLSI
- Digital maps
- Medical
- GIS
- Geophysics (Oil & Gas)
- Telecom
- Image Processing
- CAD/CAM
- Scientific visualization
Why They Use CGAL

“I recommended to the senior management that we start a policy of buying-in as much functionality as possible to reduce the quantity of code that our development team would have to maintain.

This means that we can concentrate on the application layer and concentrate on our own problem domain.”

Senior Development Engineer
& Structural Geologist

Midland Valley Exploration
“My research group JYAMITI at the Ohio State University uses CGAL because it provides an efficient and robust code for Delaunay triangulations and other primitive geometric predicates. Delaunay triangulation is the building block for many of the shape related computations that we do. [...]”

Without the robust and efficient codes of CGAL, these codes could not have been developed.”

Tamal Dey
Professor, Ohio State University
CGAL Open Source Project
Project = « Planned Undertaking »

- Institutional members make a long term commitment: Inria, MPI, Tel-Aviv U, Utrecht U, Groningen U, ETHZ, GeometryFactory, FU Berlin, Forth, U Athens

- Editorial Board
  - Steers and animates the project
  - Reviews submissions

- Development Infrastructure
  - Gforge: svn, tracker, nightly testsuite,...
  - 120p developer manual and mailing list
  - Two 1-week developer meetings per year
Contributions

- Submission of specifications of new contributions
- Review and decision by the Editorial Board

- Value for contributor
  - Integration in the CGAL community
  - Gain visibility in a mature project
  - Publication value for accepted contributions
Exact Geometric Computing
Predicates and Constructions

Predicates
- orientation
- in_circle

Constructions
- intersection
- circumcenter
Robustness Issues

• Naive use of floating-point arithmetic causes geometric algorithms to:
  – Produce [slightly] wrong output
  – Crash after invariant violation
  – Infinite loop

• There is a gap between
  – Geometry in theory
  – Geometry with floating-point arithmetic
Geometry in Theory

\[ \text{ccw}(s,q,r) \land \text{ccw}(p,s,r) \land \text{ccw}(p,q,s) \implies \text{ccw}(p,q,r) \]

Correctness proofs of algorithms rely on such theorems
The Trouble with Double

orientation(p,q,r) = sign((p_x-r_x)(q_y-r_y)-(p_y-r_y)(q_x-r_x))
Exact Geometric Computing  [Yap]

Make sure that the control flow in the implementation corresponds to the control flow with exact real arithmetic
Filtered Predicates

- Generic functor adaptor Filtered_predicate<>
  - Try the predicate instantiated with intervals
  - In case of uncertainty, evaluate the predicate with multiple precision arithmetic

- Refinements:
  - Static error analysis
  - Progressively increase precision
Filtered Constructions

Lazy number = interval and arithmetic expression tree

(3.2 + 1.5) * 13

Lazy object = approximated object and geometric operation tree

Test that may trigger an exact re-evaluation:

if ( n' < m' )

if (collinear(a',m',b'))
The User Perspective

- **Convenience Kernels**
  - `Exact_predicates_inexact_constructions_kernel`
  - `Exact_predicates_exact_constructions_kernel`
  - `Exact_predicates_exact_constructions_kernel_with_sqrt`

- **Number Types**
  - `double`, `float`
  - `CGAL::Gmpq` (rational), `Core` (algebraic)
  - `CGAL::Lazy_exact_nt<ExactNT>`

- **Kernels**
  - `CGAL::Cartesian<NT>`
  - `CGAL::Filtered_kernel<Kernel>`
  - `CGAL::Lazy_kernel<NT>`
Merits and Limitations

- Ultimate robustness inside the black box

- The time penalty is reasonable, e.g. 10% for 3D Delauny triangulation of 1M random points

Limitations of Exact Geometric Computing
- Topology preserving rounding is non-trivial
- Construction depth must be reasonable
- Cannot handle trigonometric functions
Generic Programming
template <class Key, class Less>
class set {
    Less less;

    insert(Key k)
    {
        if (less(k, treenode.key))
            insertLeft(k);
        else
            insertRight(k);
    }
};
template < class Geometry >
class Delaunay_triangulation_2 {
    Geometry::Orientation orientation;
    Geometry::In_circle in_circle;

    void insert(Geometry::Point t) {
        ...
        if(in_circle(p,q,r,t)) {...}
        ...
        if(orientation(p,q,r)){...}
    }
};
template < class Geometry, class TDS >
class Delaunay_triangulation_2 {
}

template < class Vertex, class Face >
class Triangulation_data_structure_2 {
};
Iterators

template <class Geometry>
class Delaunay_triangulation_2 {

    typedef .. Vertex_iterator;
    typedef .. Face_iterator;

    Vertex_iterator vertices_begin();
    Vertex_iterator vertices_end();

    template <class OutputIterator>
    incident_faces(Vertex_handle v, OutputIterator it);
};

std::list<Face_handle> faces;
dt.incident_faces(v, std::back_inserter(faces));
Iterators

template <class Geometry>

class Delaunay_triangulation_2 {

    template <class T>
    void insert(T begin, T end); // typeof(*begin)==Point

};

list<Kernel::Point_2> points;

Delaunay_triangulation<Kernel> dt;

dt.insert(points.begin(), points.end());
Boost Graph Library (BGL)

- Rich collection of graph algorithms: shortest paths, minimum spanning tree, flow, etc.
- Design that
  - decouples data structure from algorithm
  - links them through a thin glue layer
- BGL and CGAL
  - Provide glue layer for CGAL
  - Extension to embedded graphs inducing the notion of faces
template <typename Graph>
struct boost::graph_traits {
    typedef ... vertex_descriptor;
    typedef ... edge_descriptor;
    typedef ... vertex_iterator;
    typedef ... out_edge_iterator;
};
vertex_descriptor v, w;
edge_descriptor e;

v = source(e,G);
w = target(e,G);

std::pair<out_edge_iterator, out_edge_iterator> ipair;

ipair = out_edges(v,G);
BGL Glue Layer for CGAL

CGAL provides partial specializations:

```cpp
template <typename T>
graph_traits<Polyhedron<T>>;
```

```cpp
template <typename T>
Polyhedron<T>::Vertex
source(Polyhedron<T>::Edge);
```

Users can run:

```cpp
boost::kruskal_mst(P);
```
From A BGL Glue Layer for CGAL
To BGL Style CGAL Algorithms

BGL::Algorithm

<<Concept>>
Graph

<<Concept>>
EmbeddedGraph

boost::adjacency_list

CGAL::HDS

CGAL::Algorithm

User::Polyhedron
Summary: Overview

- Huge collection with uniform APIs
- Modular and not monolithic
- Open Source and commercial licenses

- Clear focus on geometry
- Interfaces with de facto standards/leaders: STL, Boost, GMP, Qt, blas

- Robust and fast through exact geometric computing
- Easy to integrate through generic programming
Polyhedral Surfaces

Pierre Alliez
INRIA
Outline

- Halfedge Data Structure and Polyhedron
- Euler Operators
- Customization
- Algorithms for Geometric Modelling and Geometry Processing
Halfedge Data Structure

Represented by vertices, edges, facets and an **incidence relation** on them, restricted to orientable 2-manifolds with boundary.
Polyhedron

Building blocks assembled with C++ templates
### Default Polyhedron

#### Vertex
- `Halfedge_handle halfedge()`
- `Point& point()`
- ...... ...

#### Halfedge
- `Halfedge_handle opposite()`
- `Halfedge_handle next()`
- `Halfedge_handle prev()`
- `Vertex_handle vertex()`
- `Facet_handle facet()`
- ...... ...

#### Facet
- `Halfedge_handle halfedge()`
- `Plane& plane()`
- `Normal& normal()`
- `Color& color()`
- ...... ...

---

**Diagram**

```
incident vertex

next halfedge

opposite halfedge

halfedge

previous halfedge

incident facet
```
Example

```cpp
#include <CGAL/Cartesian.h>
#include <CGAL/Polyhedron_3.h>

typedef CGAL::Simple_cartesian<double> Kernel;
typedef CGAL::Polyhedron_3<Kernel> Polyhedron;
typedef Polyhedron::Vertex_iterator Vertex_iterator;

int main()
{
    Polyhedron p;
    // ... read from file or build

    Vertex_iterator v;
    for (v = p.vertices_begin();
         v != p.vertices_end();
         ++v)
    std::cout << v->point() << std::endl;
}
Flexible Data Structure

**Vertex**
- Halfedge_handle halfedge()
- Point& point()
- ...

**Halfedge**
- Halfedge_handle opposite()
- Halfedge_handle next()
- Halfedge_handle prev()
- Vertex_handle vertex()
- Facet_handle facet()
- ...

**Facet**
- Halfedge_handle halfedge()
- Plane& plane()
- Normal& normal()
- Color& color()
- ...

![Diagram of flexible data structure](image)
Iterate over all Vertices

Vertex_iterator v;
for( v = polyhedron.vertices_begin();
    v != polyhedron.vertices_end();
    ++v )
{
    // do something with v
}
Circulate around Facet

```cpp
Halfedge_around_facet_circulator he, end;
he = end = f->facet_begin();
CGAL_For_all(he, end)
{
    // do something with he
}
```
Circulate around Vertex

```cpp
Halfedge_around_vertex_circulator he, end;
he = end = v->vertex_begin();
CGAL_For_all(he, end)
{
    // do something with he
}
```
Euler Operators

- **split_facet**
- **join_facet**
- **create_center_vertex**
- **erase_center_vertex**
- **split_vertex**
- **join_vertex** (aka edge collapse)
- **split_loop**
- **join_loop**
- **add_vertex_and_facet_to_border**
- **add_facet_to_border**
- **erase_facet**
Customization

template <class Refs>
struct MyFace : public CGAL::HalfedgeDS_face_base<Refs> {
    CGAL::Color color;
};

struct MyItems : public CGAL::Polyhedron_items_3 {
    template <class Refs, class Traits>
    struct Face_wrapper {
        typedef MyFace<Refs> Face;
    };
};

typedef CGAL::Simple_cartesian<double> Kernel;
typedef CGAL::Polyhedron_3<Kernel, MyItems> MyPolyhedron;
Algorithms

- Intersection detection
- Convex hull
- Boolean operations
- Kernel
- Parameterization
- Subdivision
- Principal component analysis
- Estimation of curvatures
- Extraction of ridges
- Simplification
Intersection Detection

Efficient algorithm for finding all intersecting pairs for large numbers of axis-aligned bounding boxes.

Generic programming:
Boxes can contain objects of any type
Example: Intersecting 3D Triangles

```cpp
#include <CGAL/Exact_predicates_inexact_constructions_kernel.h>
#include <CGAL/intersections.h>
#include <CGAL/box_intersection_d.h>

typedef CGAL::Exact_predicates_inexact_constructions_kernel Kernel;
typedef Kernel::Triangle_3 Triangle_3;
typedef std::list<Triangle_3>::iterator Iterator;
typedef CGAL::Box_intersection_d::Box_with_handle_d<double, 3, Iterator> Box;

void callback(const Box& a, const Box& b)
{
    Triangle_3 ta = *a.handle();
    Triangle_3 tb = *b.handle();
    if(CGAL::do_intersect(ta,tb)) {
        // do something
    }
}

std::list<Triangle_3> triangles(...);

std::list<Box> boxes;
Iterator it;
for(it = triangles.begin(); it != triangles.end(); ++it)
    boxes.push_back(Box((*it).bbox(),it));

CGAL::box_self_intersection_d(boxes.begin(), boxes.end(), callback);
```
Convex Hull

- From point set
- Outputs a polyhedron
#include <CGAL/convex_hull_3.h>
#include <CGAL/Exact_predicates_inexact_constructions_kernel.h>

typedef CGAL::Exact_predicates_inexact_constructions_kernel Kernel;
typedef CGAL::Convex_hull_traits_3<Kernel> Traits;
typedef Traits::Polyhedron_3 Polyhedron;
typedef Kernel::Point_3 Point;

int main()
{
    std::list<Point> points;
    // fill container...
    Polyhedron polyhedron;
    CGAL::convex_hull_3(points.begin(), points.end(), polyhedron);
    return 0;
}
Boolean Operations

Operations:
- Union
- Difference
- Intersection
- Complement

Problem: not closed, i.e., result of a Boolean operation is not necessarily a Polyhedron
Nef Polyhedron

The **3D Nef polyhedron** is a B-rep data structure which is
- closed under Boolean operations
- without enforcing regularization

**Operations:**
- Union
- Intersection
- Difference
- Complement
- Interior
- Exterior
- Boundary
- Closure
- Regularization

can evaluate a CSG-tree with halfspaces as primitives and convert it to B-rep

CGAL manual
Example

```cpp
#include <CGAL/Exact_predicates_exact_constructions_kernel.h>
#include <CGAL/Polyhedron_3.h>
#include <CGAL/Nef_polyhedron_3.h>

typedef CGAL::Exact_predicates_exact_constructions_kernel Kernel;
typedef CGAL::Polyhedron_3<Kernel> Polyhedron;
typedef CGAL::Nef_polyhedron_3<Kernel> Nef_polyhedron;

Polyhedron p1;
Polyhedron p2;

Nef_polyhedron n1(p1);
Nef_polyhedron n2(p2);

n1 -= n2; // difference

Polyhedron p3;
if(n1.is_simple())
    n1.convert_to_Polyhedron(p3);
else
    // analyze/process n1 and do something...
```

```
Example

\[ n1 \ -= \ n2; // difference \]
Kernel of a Polyhedron

- Intersection of all its interior half-spaces.
Kernel of a Polyhedron

- Intersection of all its interior half-spaces
- Uses linear programming (CGAL::QP_solver)

input polyhedron

kernel
Kernel w/ Linear Programming

Does it have a kernel?
Kernel w/ Linear Programming
Parameterization

Planar

- Conformal [Eck et al., Levy et al., Desbrun et al.]
- Mean value coordinates [Floater]
- ...
#include <CGAL/Cartesian.h>
#include <CGAL/Polyhedron_3.h>
#include <CGAL/Parameterization_polyhedron_adaptor_3.h>
#include <CGAL/parameterize.h>

Polyhedron mesh;
Mesh_adaptor_polyhedron mesh_adaptor(&mesh);
CGAL::parameterize(&mesh_adaptor);
Point_2 uv = mesh_adaptor.info(he)->uv();
Parameterization

User-provided cut graph for closed or high genus surfaces.

(conformal distortion)
Parameterization

Fixed boundary

Free boundary
Subdivision

Designed to work on CGAL polyhedron

- Catmull-Clark
- Loop
- Doo-Sabin
- Sqrt3
- ...

CGAL manual
Example

```cpp
#include <CGAL/Cartesian.h>
#include <CGAL/Polyhedron_3.h>
#include <CGAL/Subdivision_method_3.h>

typedef CGAL::Cartesian<double> Kernel;
typedef CGAL::Polyhedron_3<Kernel> Polyhedron;

using CGAL::Subdivision_method_3;

Polyhedron polyhedron;
int subdivision_depth = 3;
CatmullClark_subdivision(polyhedron, subdivision_depth);
```
Principal Component Analysis

Linear least squares fitting on sets of 3D kernel objects:

- points
- triangles for triangle meshes

CGAL manual
Example

```cpp
#include <CGAL/linear_least_squares_fitting_3.h>
// more includes and typedefs

Polyhedron mesh;

std::list<Triangle_3> triangles;
Polyhedron::Facet_iterator f;
for (f = mesh.facets_begin(); f != mesh.facets_end(); ++f) {
    const Point &a = f->halfedge()->vertex()->point();
    const Point &b = f->halfedge()->next()->vertex()->point();
    const Point &c = f->halfedge()->prev()->vertex()->point();
    triangles.push_back(Triangle_3(a, b, c));
}

Plane_3 plane;
CGAL::linear_least_squares_fitting_3( triangles.begin(),
    triangles.end(),
    plane,
    CGAL::PCA_dimension_2_tag() );
```
```cpp
#include <CGAL/linear_least_squares_fitting_3.h>
#include <boost/iterator/transform_iterator.hpp>

Polyhedron mesh;

class ToTriangle {
    Triangle_3 operator(Polyhedron::FacetIterator f) {
        return Triangle_3( f->halfedge()->vertex()->point();
                        f->halfedge()->next()->vertex()->point();
                        f->halfedge()->prev()->vertex()->point();
    }
};

Plane_3 plane;
CGAL::linear_least_squares_fitting_3(
    boost::make_transform_iterator(mesh.facets_begin(), ToTriangle()),
    boost::make_transform_iterator(mesh.facets_end(), ToTriangle()),
    plane, CGAL::PCA_dimension_2_tag());
```
Fitting Points vs Triangles

fit points

fit triangles
Estimation of Curvatures

- Estimates general differential properties (Monge form) on point sets.
- Through polynomial (d-jet) fitting
Example

```cpp
#include <CGAL/Monge_via_jet_fitting.h>
typedef CGAL::Cartesian<double> Kernel;
typedef CGAL::Monge_via_jet_fitting<Kernel> Monge_fit;
typedef Monge_fit::Monge_form Monge_form;

Monge_fit monge_fit;
Monge_form monge_form =
    monge_fit(points.begin(),
              points.end(),
              dim_fitting, dim_monge);

Vector_3 kmin = monge_form.minimal_principal_direction();
Vector_3 kmax = monge_form.maximal_principal_direction();
Vector_3 normal = monge_form.normal_direction();
```
Extraction of Ridges

Ridge: curve along which one of the principal curvatures has an extremum along its curvature line.
Simplification


BGL-style allows simplification of any model of EmbeddedGraph
**Simplification: Example**

```cpp
#include <CGAL/Simple_cartesian.h>
#include <CGAL/Polyhedron_3.h>
#include <CGAL/Surface_mesh_simplification/HalfedgeGraph_Polyhedron_3.h>
#include <CGAL/Surface_mesh_simplification/edge_collapse.h>
#include <CGAL/Surface_mesh_simplification/Policies/Edge_collapse/Count_stop_predicate.h>

typedef CGAL::Simple_cartesian<double> Kernel;
typedef CGAL::Polyhedron_3<Kernel> Mesh;
namespace SMS = CGAL::Surface_mesh_simplification ;

Mesh mesh;
SMS::Count_stop_predicate<Mesh> stop(1000); // target # edges
SMS::edge_collapse(mesh, stop, 
                 CGAL::vertex_index_map( boost::get(CGAL::vertex_external_index, mesh) )
                 .edge_index_map( boost::get(CGAL::edge_external_index, mesh) ));
```

[CGAL manual]
[boost::property_map]
[bgl_named_params]
Summary and Outlook

- The halfedge data structure and the polyhedron are highly flexible
- CGAL provides algorithms for geometric modeling and geometry processing
- Polyhedral surface as output of surface mesh generation algorithms (Part IV)
Under Development

- BGL-ization of existing CGAL algorithms
Under Development

- BGL-ization of existing CGAL algorithms
- Remeshing
Questions?
Arrangements

Efi Fogel
Tel Aviv University
Outline

- Arrangements
- Algorithms based on Arrangements
  - Boolean Set Operations
  - Minkowski Sums and Polygon Offset
  - Envelopes
- Arrangements on Surfaces
Arrangement Definition

Given a collection of curves on a surface, the **arrangement** is the partition of the surface into **vertices**, **edges** and **faces** induced by the curves.

- An arrangement of circles in the plane
- An arrangement of lines in the plane
- An arrangement of geodesic arcs on the sphere
```cpp
#include <CGAL/Exact_predicates_exact_constructions_kernel.h>
#include <CGAL/Arr_segment_traits_2.h>
#include <CGAL/Arrangement_2.h>

typedef CGAL::Exact_predicates_exact_constructions_kernel Kernel;
typedef CGAL::Arr_segment_traits_2<Kernel> Traits;
typedef Traits::Point_2 Point;
typedef Traits::X_monotone_curve_2 Segment;
typedef CGAL::Arrangement_2<Traits> Arrangement;

int main() {
    Point p1(0, 0), p2(1, 0), p3(0, 1);
    Segment cv[3] = { Segment(p1,p2), Segment(p2,p3), Segment(p3,p1) };
    Arrangement arr;
    CGAL::insert(arr, cv, cv+3);
    return (arr.is_valid()) ? 0 : -1;
}
```

This code defines a simple arrangement of three line segments, forming a triangle with vertices at points p1, p2, and p3. The `main` function creates the arrangement, inserts the line segments, and checks its validity.
CGAL::Arrangement_2

- Constructs, maintains, modifies, traverses, queries, and presents subdivisions of the plane
- Robust and exact
  - All inputs are handled correctly (including degenerate)
  - Exact number types are used to achieve exact results
- Efficient
- Generic
  - Easy to interface, extend, and adapt
  - Notification mechanism for change propagation
- Modular
  - Geometric and topological aspects are separated
Geometric Traits

- Define the family of curves
- Aggregate geometric types and operations over the types

<table>
<thead>
<tr>
<th>Compare two points</th>
</tr>
</thead>
<tbody>
<tr>
<td>Determine the relative position of a point and an $x$-monotone curve</td>
</tr>
<tr>
<td>Determine the relative position of two $x$-monotone curves to the left (right) of a point</td>
</tr>
<tr>
<td>Subdivide a curve into $x$-monotone curves</td>
</tr>
<tr>
<td>Find all intersections of two $x$-monotone curves</td>
</tr>
</tbody>
</table>
**Geometric Traits**

- Define the family of curves
- Aggregate geometric types and operations over the types
- Compare two points
- Find all intersections of two $x$-monotone curves
- Subdivide a curve into $x$-monotone curves
- Determine the relative position of two $x$-monotone curves to the left (right) of a point
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**Geometric Traits**

- Define the family of curves
- Aggregate geometric types and operations over the types

<p>| | |</p>
<table>
<thead>
<tr>
<th></th>
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<tbody>
<tr>
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<tr>
<td></td>
<td></td>
</tr>
</tbody>
</table>

- Compare two points

- Determine the relative position of a point and an $x$-monotone curve

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  - Subdivide a curve into $x$-monotone curves
  - Find all intersections of two $x$-monotone curves
### Arrangement Traits Classes

<table>
<thead>
<tr>
<th>Curve Family</th>
<th>Degree</th>
<th>Surface</th>
<th>Boundness</th>
<th>Arithmetic</th>
<th>Attribute</th>
</tr>
</thead>
<tbody>
<tr>
<td>linear segments</td>
<td>1</td>
<td>plane</td>
<td>bounded</td>
<td>rational</td>
<td>caching</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td>noncaching</td>
</tr>
<tr>
<td>linear segments, rays, lines</td>
<td>1</td>
<td>plane</td>
<td>unbounded</td>
<td>rational</td>
<td></td>
</tr>
<tr>
<td></td>
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<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>piecewise linear curves</td>
<td>$\infty$</td>
<td>plane</td>
<td>bounded</td>
<td>rational</td>
<td>caching</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td>noncaching</td>
</tr>
<tr>
<td>circular arcs, linear segments</td>
<td>$\leq 2$</td>
<td>plane</td>
<td>bounded</td>
<td>rational</td>
<td></td>
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<tr>
<td></td>
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<td></td>
</tr>
<tr>
<td>algebraic curves</td>
<td>$\leq 2$</td>
<td>plane</td>
<td>Bounded</td>
<td>algebraic</td>
<td>CORE</td>
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<td></td>
<td></td>
<td></td>
<td>unbounded</td>
<td></td>
<td>CKvA_2</td>
</tr>
<tr>
<td>quadric projections</td>
<td>$\leq 2$</td>
<td>plane</td>
<td>unbounded</td>
<td>algebraic</td>
<td></td>
</tr>
<tr>
<td>algebraic curves</td>
<td>$\leq 3$</td>
<td>plane</td>
<td>unbounded</td>
<td>algebraic</td>
<td></td>
</tr>
<tr>
<td>algebraic curves</td>
<td>$\leq n$</td>
<td>plane</td>
<td>unbounded</td>
<td>algebraic</td>
<td></td>
</tr>
<tr>
<td>planar Bézier curves</td>
<td>$\leq n$</td>
<td>plane</td>
<td>unbounded</td>
<td>algebraic</td>
<td></td>
</tr>
<tr>
<td>univariate polynomials</td>
<td>$\leq n$</td>
<td>plane</td>
<td>unbounded</td>
<td>algebraic</td>
<td>RS</td>
</tr>
<tr>
<td>rational function arcs</td>
<td>$\leq n$</td>
<td>plane</td>
<td>unbounded</td>
<td>algebraic</td>
<td></td>
</tr>
<tr>
<td>geodesic arcs on sphere</td>
<td>$\leq 2$</td>
<td>sphere</td>
<td>bounded</td>
<td>rational</td>
<td></td>
</tr>
<tr>
<td>quadric intersection arcs</td>
<td>$\leq 2$</td>
<td>quadric</td>
<td>unbounded</td>
<td>algebraic</td>
<td></td>
</tr>
<tr>
<td>dupin cyclide intersection. arcs</td>
<td>$\leq 2$</td>
<td>dupin cyclides</td>
<td>bounded</td>
<td>algebraic</td>
<td></td>
</tr>
</tbody>
</table>
template <typename Traits, typename Dcel = Arr_default_dcel<Traits_> >
class Arrangement_2 {

...
};

enum Color {BLUE, RED, WHITE};

typedef CGAL::Arr_segment_traits_2<Kernel> Traits;
typedef Traits::Point_2 Point;
typedef Traits::X_monotone_curve_2 Segment;
typedef CGAL::Arr_extended_dcel<Traits, Color, Color, Color> Dcel;
typedef CGAL::Arrangement_2<Traits, Dcel> Arrangement;
Point Location

Given a subdivision $A$ of the space into cells and a query point $q$, find the cell of $A$ containing $q$

- Fast query processing
- Reasonably fast preprocessing
- Small space data structure

<table>
<thead>
<tr>
<th></th>
<th>Naive</th>
<th>Walk</th>
<th>RIC</th>
<th>Landmarks</th>
</tr>
</thead>
<tbody>
<tr>
<td>Preprocessing time</td>
<td>none</td>
<td>none</td>
<td>$O(n \log n)$</td>
<td>$O(k \log k)$</td>
</tr>
<tr>
<td>Memory space</td>
<td>none</td>
<td>none</td>
<td>$O(n)$</td>
<td>$O(k)$</td>
</tr>
<tr>
<td>Query time</td>
<td>bad</td>
<td>reasonable</td>
<td>good</td>
<td>good</td>
</tr>
<tr>
<td>Applicability</td>
<td>all</td>
<td>limited</td>
<td>limited</td>
<td>limited</td>
</tr>
</tbody>
</table>

**Walk** — Walk along a line

**RIC** — Random Incremental Construction based on trapezoidal decomposition

$k$ — number of landmarks
typedef CGAL::Arr_naive_point_location<Arrangement_2> Naive_pl;
typedef CGAL::Arr_landmarks_point_location<Arrangement_2> Landmarks_pl;

int main() {
    Arrangement arr;
    construct_arr(arr);
    Point p(1, 4);

    Naive_pl naive_pl(arr);    // Associate arrangement to naïve point location
    CGAL::Object obj1 = naive_pl.locate(p);

    Landmarks landmarks_pl;
    landmarks_pl.attach(arr);   // Attach landmarks point location to arrangement
    CGAL::Object obj2 = landmarks_pl.locate(p);

    return (obj1 == obj2) ? 0 : -1;
}
Boolean Set Operations

For two point sets $P$ and $Q$ and a point $r$:

<table>
<thead>
<tr>
<th>Operation</th>
<th>Expression</th>
</tr>
</thead>
<tbody>
<tr>
<td>Complement</td>
<td>$R = \overline{P}$</td>
</tr>
<tr>
<td>Union</td>
<td>$R = P \cup Q$</td>
</tr>
<tr>
<td>Intersection</td>
<td>$R = P \cap Q$</td>
</tr>
<tr>
<td>Difference</td>
<td>$R = P \setminus Q$</td>
</tr>
<tr>
<td>Symmetric Difference</td>
<td>$R = (P \setminus Q) \cup (Q \setminus P)$</td>
</tr>
</tbody>
</table>

Intersection predicate: $P \cap Q \neq \emptyset$ Overlapping cell(s) are not explicitly computed

Containment predicate: $r \in P$

Interior, Boundary, Closure

Regularization: $R = \text{closure}(\text{interior}(P))$
int main() {
    Polygon p, q;
    p.push_back(Point(0, 0)); p.push_back(Point(2, 0));
    p.push_back(Point(1, 1)); p.push_back(Point(0, 2));
    q.push_back(Point(1, 1)); q.push_back(Point(3, 1));
    q.push_back(Point(2, 2)); q.push_back(Point(1, 3));

    Polygon_with_holes comp_p, comp_q;
    CGAL::complement(p, comp_p);
    CGAL::complement(q, comp_q);

    Polygon_with_holes a;
    CGAL::join(comp_p, comp_q, a);

    std::list<Polygon_with_holes> l1, l2;
    CGAL::complement(a, std::back_inserter(l1));
    CGAL::intersection(p, q, l2);

    return std::compare(l1, l2);
}
CGAL::Boolean_set_operation_2

- Supports
  - regularized Boolean set-operations
  - intersection predicates
  - point containment predicates
- Operands and results are regularized point sets bounded by $x$-monotone curves referred to as general polygons
  - General polygons may have holes
- Extremely efficient aggregated operations
- Based on the `Arrangement_2` and `Polygon_2` packages
Minkowski Sum in $\mathbb{R}^d$

$P$ and $Q$ are 2 polytopes in $\mathbb{R}^d$

$$P \oplus Q = \{ p + q \mid p \in P, q \in Q \}$$

$P \cap Q \neq \emptyset \iff \text{Origin} \in P \oplus (-Q)$

collision
CGAL::Minkowski_sum_2

- Based on the Arrangement_2, Polygon_2, and Partition_2 packages
- Supports Minkowski sums of two simple polygons
  - Implemented using either decomposition or convolution
  - Exact
- Interoperable with Boolean_set_operations_2, e.g., Compute the union of offset polygons
- Supports Minkowski sums of a simple polygon and a disc (polygon offsetting)
  - Offers either an exact computation or a conservative approximation scheme
  - Disk radius can be negative (inner offset)
Motion Planning

- The input robot and the obstacle are non-convex
- Exploits the convolution method
- The output sum contains four holes, isolated points, and antennas
Envelopes in $\mathbb{R}^3$

The **lower envelope** of a set of $xy$-monotone surfaces $S = \{S_1, S_2, \ldots, S_n\}$, is the point-wise minimum of all surfaces.

The **minimization diagram** of $S$ is an arrangement.

- The identity of the surfaces that induce the lower envelope over a specific cell (vertex, edge, face) of the arrangement is the same.
```cpp
#include <CGAL/Exact_predicates_exact_constructions_kernel.h>
#include <CGAL/Env_triangle_traits_3.h>
#include <CGAL/Env_surface_data_traits_3.h>
#include <CGAL/envelope_3.h>

typedef CGAL::Exact_predicates_exact_constructions_kernel Kernel;
typedef CGAL::Env_triangle_traits_3<Kernel>               Traits;
typedef Traits::Point_3                                   Point;
typedef Traits::Surface_3                                 Tri;
enum Color {RED, BLUE};
typedef CGAL::Env_surface_data_traits_3<Traits, Color>    Data_traits;
typedef Data_traits::Surface_3                            Dtri;
typedef CGAL::Envelope_diagram_2<Data_traits>             Arrangement;

int main() {
    Point p1(0,0,1), p2(0,6,1), p3(5,3,5), p4(6,0,1), p5(6,6,1), p6(1,3,5);
    Tri tris[] = {Dtri(Tri(p1,p2,p3), RED), Dtri(Tri(p4,p5,p6), BLUE)};
    Arrangement arr;
    CGAL::lower_envelope_3(tris, tris+2, arr);  
    return (arr.is_valid()) ? 0 : -1;
}
```
CGAL::Envelope_3

- Constructs lower and upper envelopes of surfaces

<table>
<thead>
<tr>
<th>Surface Family</th>
<th>Class Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>triangles</td>
<td>Env_triangle_traits_3</td>
</tr>
<tr>
<td>spheres</td>
<td>Env_sphere_traits_3</td>
</tr>
<tr>
<td>planes and half planes</td>
<td>Env_plane_traits_3</td>
</tr>
<tr>
<td>quadrics</td>
<td>Env_quadric_traits_3</td>
</tr>
</tbody>
</table>

- Based on the **Arrangement_2** package

- Exploits
  - Overlay computation (using the sweep line framework)
  - Isolated points
  - Zone traversal
Lower Envelopes
Voronoi Diagrams via Envelopes

- Computed as lower envelopes of planes
- Represented as planar arrangements of unbounded curves

points along a line segment
points on a grid inside a square
points on a circle
points on a square boundary
Arrangements on Surfaces in $\mathbb{R}^3$

A parametric surface $S$ of two parameters is a surface defined by parametric equations involving two parameters $u$ and $v$:

$$f_S(u, v) = (x(u, v), y(u, v), z(u, v))$$

Thus, $f_S : P \rightarrow \mathbb{R}^3$ and $S = f_S(P)$, where $P$ is a continuous and simply connected two-dimensional parameter space.

We deal with orientable parametric surfaces.
Minkowski-Sums of Polytopes

- The Gaussian map of a polytope is the decomposition of $S^2$ into maximal connected regions so that the extremal point is the same for all directions within one region.
- The overlay of the Gaussian maps of two polytopes $P$ and $Q$ is the Gaussian map of the Minkowski sum of $P$ and $Q$. 
Voronoi Diagrams on the Sphere

- All algorithms supported by the `Arrangement_2` package can also be used on the sphere.
- We compute lower envelopes defined over the sphere.
- We can compute Voronoi diagrams on the sphere, the edges of which are geodesic arcs.
Arrangements on the Sphere

- The overlay of
  - An arrangement on the sphere induced by
    - the continents and some of the islands on earth
    - 5 cities
      - New Orleans
      - Los Angeles
      - San Antonio
      - San Diego
      - Boston
    - Voronoi diagram of the cities

- Diagrams, envelopes, etc. are represented as arrangements
  - Can be passed as input to consecutive operations

The sphere is oriented such that Cambridge is at the center
Video

- Arrangements of Geodesic Arcs on the Sphere
- Was presented at the 24th ACM Symposium on Computational Geometry, College Park, Maryland, July 2008
Summary

- Arrangements are versatile tools
- Arrangements are used as foundation for higher level geometric data structures
- Arrangements are not bound to the plane
Triangulations and Mesh Generation

Andreas Fabri
GeometryFactory

Pierre Alliez
INRIA
Outline

• **2D**
  • From triangulations to quality meshes
  • Related components

• **3D**
  • Triangulations
  • Mesh generation
    • Key concepts
    • Surfaces
    • Volumes
    • Next release and work in progress

• **Summary**
2D
From Triangulations to Quality Meshes
Delaunay Triangulation

- A triangulation is a Delaunay triangulation, if the circumscribing circle of any facet of the triangulation contains no vertex in its interior.
```cpp
#include <CGAL/Exact_predicates_inexact_constructions_kernel.h>
#include <CGAL/Delaunay_triangulation_2.h>

typedef CGAL::Exact_predicates_inexact_constructions_kernel Kernel;
typedef Kernel::Point_2 Point;

typedef CGAL::Delaunay_triangulation_2<Kernel> Delaunay;
typedef Delaunay::Vertex_handle Vertex_handle;

int main()
{
    Delaunay dt;

    dt.insert( std::istream_iterator<Point>(std::cin),
               std::istream_iterator<Point>() );

    Vertex_handle v = dt.nearest_vertex(Point(0.0,0.0));

    std::cout << "Nearest vertex to origin: " << v->point() << std::endl;

    return 0;
}
```
Adding Constraints
Include the necessary headers:

```cpp
#include <CGAL/Exact_predicates_inexact_constructions_kernel.h>
#include <CGAL/Constrained_Delaunay_triangulation_2.h>
```

Create a Kernel type:

```cpp
typedef CGAL::Exact_predicates_inexact_constructions_kernel Kernel;
```

Create a Point type:

```cpp
typedef Kernel::Point_2 Point;
```

Create a Constrained Delaunay Triangulation type:

```cpp
typedef CGAL::Constrained_Delaunay_triangulation_2<Kernel> CDT;
```

Create a Vertex Handle type:

```cpp
typedef CDT::Vertex_handle Vertex_handle;
```

Define the main function:

```cpp
int main()
{
    CDT cdt;

    // from points
    cdt.insert_constraint(Point(0.0,0.0), Point(1.0,0.0));

    // from vertices
    Vertex_handle v1 = cdt.insert(Point(2.0,3.0));
    Vertex_handle v2 = cdt.insert(Point(4.0,5.0));
    cdt.insert_constraint(v1,v2);

    return 0;
}
```
Conforming Delaunay
Code Example

```cpp
#include <CGAL/Triangulation_conformer_2.h>

// constrained Delaunay triangulation
CDT cdt;
... // insert points & constraints

CGAL::make_conforming_Delaunay_2(cdt);
```
Delaunay Meshing
#include <CGAL/Exact_predicates_inexact_constructions_kernel.h>
#include <CGAL/Constrained_Delaunay_triangulation_2.h>
#include <CGAL/Delaunay_mesher_2.h>

typedef CGAL::Exact_predicates_inexact_constructions_kernel Kernel;

typedef CGAL::Constrained_Delaunay_triangulation_2<Kernel> CDT;

int main()
{
    CDT cdt;
    ... // insert points and constraints
    CGAL::refine_Delaunay_mesh_2(cdt);
    return 0;
}
Background
Delaunay Edge

An edge is said to be a Delaunay edge, if it is inscribed in an empty circle.
Gabriel Edge

An edge is said to be a **Gabriel edge**, if its diametral circle is empty.
A constrained Delaunay triangulation is a conforming Delaunay triangulation, if every constrained edge is a Delaunay edge.
A constrained Delaunay triangulation is a conforming Gabriel triangulation, if every constrained edge is a Gabriel edge.
Steiner Vertices

Any constrained Delaunay triangulation can be refined into a conforming Delaunay or Gabriel triangulation by adding Steiner vertices.
Delaunay Refinement

**Rule #1:** break bad elements by inserting circumcenters (Voronoi vertices)
- “bad” in terms of size or shape (too big or skinny)

Picture taken from [Shewchuk]
Delaunay Refinement

Rule #2: Midpoint vertex insertion
A constrained segment is said to be *encroached*, if there is a vertex inside its diametral circle.

Picture taken from [Shewchuk]
Delaunay Refinement

Encroached subsegments have priority over skinny triangles

Picture taken from [Shewchuk]
API
Rich API

- Traversal
Rich API

- Traversal
- Localization

Online manual
Rich API

- Traversal
- Localization
- Dynamic: insertion & removal
Rich API

- Traversal
- Localization
- Dynamic: insertion & removal
- Parameters for mesh generation
Parameters for Mesh Generation

- **Shape**
  - Lower bound on triangle angles

Input PLSG

5 deg

20.7 deg

Online manual
Parameters for Mesh Generation

- **Shape**
  - Lower bound on triangle angles

- **Size**
  - No constraint
  - Uniform sizing
  - Sizing function
Sizing Parameter

No constraint  Uniform  Sizing function
Example Code

```cpp
#include <CGAL/Exact_predicates_inexact_constructions_kernel.h>
#include <CGAL/Constrained_Delaunay_triangulation_2.h>
#include <CGAL/Delaunay_mesher_2.h>
#include <CGAL/Delaunay_mesh_size_criteria_2.h>

typedef CGAL::Exact_predicates_inexact_constructions_kernel Kernel;
typedef CGAL::Constrained_Delaunay_triangulation_2<Kernel> CDT;
typedef CGAL::Delaunay_mesh_size_criteria_2<CDT> Criteria;
typedef CGAL::Delaunay_mesher_2<CDT, Criteria> Meshing_engine;

int main()
{
    CDT cdt;
    Meshing_engine engine(cdt);
    engine.refine_mesh();
    engine.set_criteria(Criteria(0.125, 0.5)); // min 20.6 deg
        // 0.5 for sizing
    engine.refine_mesh(); // refine once more, etc.
    return 0;
}
```
Parameters for Mesh Generation

- **Shape**
  - Lower bound on triangle angles

- **Size**
  - No constraint
  - Uniform sizing
  - Sizing function

- **Seeds**
  - Exclude/include components
Performances

Kilimandjaro: elevation contour lines (38K segments)

Online demo
Performances

Refinement: 15K vertices/s
Related Components

• Voronoi diagram
Related Components

• Voronoi diagram
• Elevation (through traits class)

Courtesy IPF, Vienna University of Technology & Inpho GmbH
Related Components

- Voronoi diagram
- Elevation
- Interpolation (natural neighbors)
Related Components

- Voronoi diagram
- Elevation
- Interpolation
- Placement of streamlines
3D
3D Triangulations

- Delaunay
- Regular
- Rich API
- Fully dynamic
- 1M points in 16s

Online manual

Tetrahedron
Mesh Generation

Key concepts:
- Delaunay filtering
- Delaunay refinement
Delaunay Filtering

Dual Voronoi edge

Voronoi edge $\cap$ surface $S$

Delaunay triangulation restricted to surface $S$
Delaunay Refinement

Steiner point

**Bad** facet = big or badly shaped or large approximation error
Surface Mesh Generation Algorithm

repeat
{
    pick bad facet \( \mathbf{f} \)
    insert furthest \((\text{dual}(\mathbf{f}) \cap \mathbf{S})\) in Delaunay triangulation
    update Delaunay triangulation restricted to \( \mathbf{S} \)
}
until all facets are good
Isosurface from 3D Grey Level Image

```cpp
#include <CGAL/Surface_mesh_default_triangulation_3.h>
#include <CGAL/Complex_2_in_triangulation_3.h>
#include <CGAL/make_surface_mesh.h>
#include <CGAL/Gray_level_image_3.h>
#include <CGAL/Implicit_surface_3.h>

typedef CGAL::Surface_mesh_default_triangulation_3 Tr;
typedef CGAL::Complex_2_in_triangulation_3<Tr> C2t3;
typedef CGAL::Implicit_surface_3<Kernel, Gray_level_image> Surface_3;

Tr tr; // 3D-Delaunay triangulation
C2t3 c2t3 (tr); // 2D-complex in 3D-Delaunay triangulation
Gray_level_image image("data/brain",128);
Surface_3 surface(image, bounding_sphere, 1e-2);

// Criteria: min triangle angles, size, approximation error,
CGAL::Surface_mesh_default_criteria_3<Tr> criteria(30.,5.,5.);
CGAL::make_surface_mesh(c2t3, surface, criteria, CGAL::Manifold_tag());
```
Output Mesh

Triangle surface mesh approximating $S$

input
Output Mesh Properties

• Well shaped triangles
  – Lower bound on triangle angles
• Homeomorphic to input surface
• Manifold
  – not only combinatorially, i.e., no self-intersection
• Faithful Approximation of input surface
  – Hausdorff distance
  – Normals
Delaunay Refinement vs Marching Cubes

Delaunay refinement

Marching cubes in octree
Surface Remeshing

Input is a polyhedral surface

(requires efficient data structures for intersection computations)
Parameters

- Shape of triangles
  - lower bound on triangle angles
- Size
Parameters

• Shape of triangles
  – lower bound on triangle angles

• Size
  – No constraint
Parameters

• Shape of triangles
  – lower bound on triangle angles

• Size
  – No constraint
  – Uniform sizing
Parameters

- **Shape of triangles**
  - lower bound on triangle angles

- **Size**
  - No constraint
  - Uniform sizing
  - Sizing function
Parameters

• Shape of triangles
  – lower bound on triangle angles

• Size
  – No constraint
  – Uniform sizing
  – Sizing function

• Approximation error
Uniform vs Adapted
Mesh Generation Framework

Refinement
while(simplex is bad)
refine(simplex)

Oracle
3D image
Surface mesh
Implicit function
...

Constraints
Size
Shape
Approximation
...

...
A Versatile Framework

- 3D grey level images
- 3D multi-domain images
- Implicit function: $f(x, y, z) = \text{constant}$
- Surface mesh (remeshing)
- Point set (surface reconstruction)
- Anything which provides intersections
Next Release
Volume Mesh Generation
More Delaunay Filtering

Delaunay triangulation restricted to domain $\Omega$

Dual Voronoi vertex inside domain $\Omega$ ("oracle")

tetrahedron

circumsphere
Delaunay Refinement

Steiner point

Bad tetrahedron = big or badly shaped
Volume Mesh Generation Algorithm

repeat
{
    pick bad simplex
    if (Steiner point encroaches a facet)
        refine facet
    else
        refine simplex
        update Delaunay triangulation restricted to domain
}
until all simplices are good
Exude slivers
Tetrahedron Zoo

sliver
Sliver Exudation [Edelsbrunner-Guoy]

- Delaunay triangulation turned into a regular triangulation with null weights.
- Small increase of weights triggers edge-facets flips to remove slivers.
Sliver Exudation Process

- **Try** improving all tetrahedra with an aspect ratio lower than a given bound
- **Never flips a boundary facet**
3D Mesh from Multi-Domain Images

Tr tr;  // 3D Delaunay triangulation
C2t3 c2t3(tr);  // 2D complex in 3D-Delaunay triangulation

Image_3 image("segmented_visible_human");
Mesh_traits mesh_traits(image);
Facets_criteria facets_criteria(5, 1);  // facet sizing
    // approximation error
Tets_criteria tets_criteria(5);  // tet sizing

// 0.5 = radius-radius ratio upper bound for sliver exudation
CGAL::make_mesh_3_for_multivolumes(c2t3, mesh_traits,
    facets_criteria,
    tets_criteria, 0.5);
Output Volume Mesh

Visible human
Work in Progress
Piecewise Smooth Surfaces
Input: Piecewise smooth complex corner
More Delaunay Filtering

<table>
<thead>
<tr>
<th>primitive</th>
<th>dual of</th>
<th>test</th>
<th>against</th>
</tr>
</thead>
<tbody>
<tr>
<td>Voronoi vertex</td>
<td>tetrahedron</td>
<td>inside</td>
<td>domain</td>
</tr>
<tr>
<td>Voronoi edge</td>
<td>facet</td>
<td>intersect</td>
<td>domain boundary</td>
</tr>
<tr>
<td>Voronoi face</td>
<td>edge</td>
<td>intersect</td>
<td>crease</td>
</tr>
</tbody>
</table>
Delaunay Refinement

- Steiner points
Summary
Summary

- From triangulation to quality meshes
- Mesh generation:
  - 2D: Preserves constraints exactly.
Summary

- From triangulation to quality meshes
- Mesh generation:
  - 2D: Preserves constraints exactly.
  - 3D: Interpolates boundary and sharp creases.
Summary

• From triangulation to quality meshes
• Mesh generation:
  – 2D: Preserves constraints exactly.
  – 3D:
    • Interpolates boundary and sharp creases.
    • Versatile through oracle-based design
See Also

DelPSC software
(based on CGAL)
[Dey-Levine]

Skin surfaces

Online manual